

Market Design

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Outline of the class

Lecture 1: Introduction

Lecture 2: Auction theory and design

Lecture 3: Common-value and Multi-unit auctions

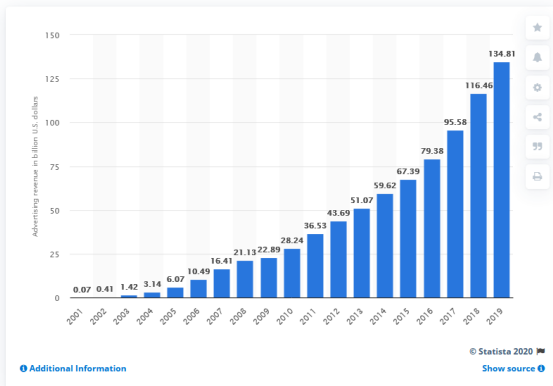
Lecture 4: Multi-item auctions and matching, sponsored-search auctions, spectrum auctions, package auctions

Sponsored Search Auctions

- Search advertising is a huge auction market.
- Why would you use auctions in this setting?
 - ▶ Difficult to set so many prices (tens of millions of keywords).
 - ▶ Demand and especially supply might be changing.
 - ▶ Retain some price-setting ability via auction design.

Advertising revenue of Google from 2001 to 2019

(in billion U.S. dollars)



Keyword Auctions

- Advertisers submit bids for keywords
 - ▶ E.g. offer a dollar payment per click.
- Separate auction for every query
 - ▶ Positions awarded in order of bid (more on this later).
 - ▶ Google uses “generalized second price” auction format (GSP).
 - ▶ Advertisers pay bid of the advertiser in the position below.
- Some important features
 - ▶ Value is created by getting a good match of ad to searcher.
 - ▶ Multiple positions, but advertisers submit only a single bid.

A Brief History

- Early 1990s: Websites sell advertising space on a “per-eyeball” basis, with contracts negotiated by salespeople; similar to print or television.
- Mid 1990s: Overture Services, Inc. (formerly GoTo.com) allows advertisers to bid for keywords, offering to pay per click. Yahoo! and others adopt this approach, charging advertisers their bids.
- 2000s: Google and Overture modify keyword auction to have advertisers pay minimum amount necessary to maintain their position (GSP).
- Late 2000s: Auction design becomes more sophisticated; auctions used to allocate advertising on many webpages (e.g. Facebook), not just search.

Example

- Two positions receive 200 and 100 clicks per day (cpd).
- 3 bidders with per click values (pcv) \$10, \$4, \$2. One position per bidder.

	Top (200 cpd)	2nd (100 cpd)
Bidder 1	2000	1000
Bidder 2	800	400
Bidder 3	400	200

- Efficient allocation creates value \$2400
- Can we find per-click prices for the two positions that clear this market?

Market Clearing Prices

- Solve for lowest market clearing prices

	Top (200 cpd)	2nd (100 cpd)
Bidder 1	2000	1000
Bidder 2	800	400
Bidder 3	400	200

- Lowest market clearing prices: $P_1 = 600$ and $P_2 = 200$
 - ▶ Bidder 1 prefers top position.
 - ▶ Bidder 2 prefers 2nd position.
 - ▶ Bidder 3 prefers to stay out.
- Corresponding per-click prices: $p_1 = 3$ and $p_2 = 2$.
- These prices achieve the efficient allocation and generate revenue of \$800.

Market Clearing Prices

- Solve for all market clearing prices

	Top (200 cpd)	2nd (100 cpd)
Bidder 1 (10 pcv)	2000	1000
Bidder 2 (4 pcv)	800	400
Bidder 3 (2 pcv)	400	200

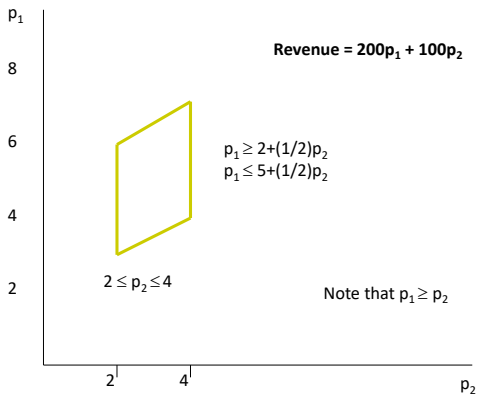
- Bidder 3 prefers nothing: $p_1 \geq 2$ and $p_2 \geq 2$
- Bidder 2 prefers 2nd: $p_2 \leq 4$ and

$$200(4 - p_1) \leq 100(4 - p_2) \Rightarrow p_1 \geq 2 + \frac{1}{2}p_2$$

- Bidder 1 prefers top: $p_1 \leq 10$ (redundant) and

$$200(10 - p_1) > 100(10 - p_2) \Rightarrow p_1 < 5 + \frac{1}{2}p_2$$

Market Clearing Prices



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Price premium for more clicks

- At market clearing prices, bidder k prefers position k to $k - 1$:

$$x_k(v_k - p_k) \geq x_{k-1}(v_k - p_{k-1}) \Rightarrow p_{k-1} \geq \underbrace{v_k}_{\geq p_k} \underbrace{\left(1 - \frac{x_k}{x_{k-1}}\right)}_{>0} + \frac{x_k}{x_{k-1}} p_k$$
$$\Rightarrow p_{k-1} \geq p_k$$

- Positions with more clicks get higher per click prices.

Finding Market Clearing Prices

Find lowest prices:

- Set $p_K = v_{K+1}$ so bidder $K + 1$ is indifferent between K and nothing.
- Set p_{k-1} so that bidder k is indifferent between k and $k - 1$:

$$x_{k-1}(v_k - p_{k-1}) = x_k(v_k - p_k)$$

Find highest prices:

- Set $p_K = v_K$ so bidder K is indifferent between K and nothing.
- Set p_{k-1} so that bidder $k - 1$ is indifferent between k and $k - 1$:

$$x_{k-1}(v_{k-1} - p_{k-1}) = x_k(v_{k-1} - p_k)$$

Sponsored Search Auction

- Can we design an auction to find market clearing prices?
- Ideally want to use the structure of the problem to design a simple auction. We will consider several options.

Overture “Pay-as-Bid” Auction:

- Each bidder submits a single bid (in \$ per click)
- Top bid gets position 1, second bid position 2, etc.
- Bidders pay their bid for each click they get.

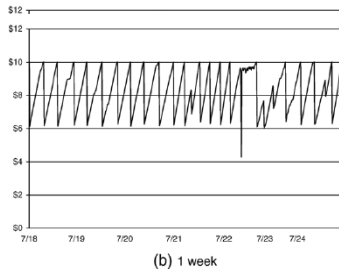
Pay-as-Bid Example

	Top (200 cpd)	2nd (100 cpd)
Bidder 1 (10 pcv)	2000	1000
Bidder 2 (4 pcv)	800	400
Bidder 3 (2 pcv)	400	200

- Bidder 3 will offer up to \$2 per click.
- Bidder 2 wants to bid \$2.01 to get second slot.
- Bidder 1 wants to bid \$2.02 to get first slot.
- But then bidder 2 wants to bid \$2.03 and get first slot...

Pay as bid auction is unstable

Overture Bid Pattern



Edelman and Ostrovsky (2006): “sawtooth” pattern caused by auto-bidding programs.

Generalized Second Price Auction:

- Bidders submit bids (in \$ per click)
- Top bid gets slot 1, second bid gets slot 2, etc.
- Each bidder pays the bid of the bidder below him.
- Seems intuitively like a more stable auction.
- Do the bidders want to bid truthfully?

Truthful Bidding

It is not a dominant strategy to bid truthfully

Example:

- Two positions: $x_1 = 200$ and $x_2 = 100$.
- Consider a bidder with value 10
- Suppose competing bids are 8 and 4
 - ▶ Bid 10 $\Rightarrow 200(10 - 8) = 400$
 - ▶ Bid 7 $\Rightarrow 100(10 - 4) = 600$

GSP Auction: Example

	Top (200 cpd)	2nd (100 cpd)
Bidder 1 (10 pcv)	2000	1000
Bidder 2 (4 pcv)	800	400
Bidder 3 (2 pcv)	400	200

- In this example, truthful bidding (10,4,2) is a Nash equilibrium (not true in general):
 - ▶ Bidder 3 would have to pay 4 to get slot 2 $\Rightarrow -200 < 0$
 - ▶ Bidder 2 gets slot 2 and profit $200 > 0$, if bid 10 to get slot 1, profit $-1200 < 200$
 - ▶ Bidder 1 has profit 1200, if bid 4 to get slot 2 $\Rightarrow (10 - 2)100 < 1200$.

GSP Auction: Example

	Top (200 cpd)	2nd (100 cpd)
Bidder 1 (10 pcv)	2000	1000
Bidder 2 (4 pcv)	800	400
Bidder 3 (2 pcv)	400	200

- This is not the only NE however: $(10,7,4)$ is also a NE
 - ▶ Bidder 3 cannot do better than 0 by changing bid.
 - ▶ Bidder 2 is making 0 profit, but would get -600 if changed his bid to 10.
 - ▶ Bidder 1 makes 600 profit, and would get $(10 - 4)100 = 600$ if changed bid to 7 to get slot 2.
- This equilibrium achieves the highest market clearing prices, i.e. highest revenue.

Assumption: $N > K$.

Theorem. For every set of market clearing prices p_1, \dots, p_K , there exists a NE of the GSP auction that achieves these prices.

Proof.

- Bidder 1 bids his value, bidder 2 bids p_1, \dots , bidder $K + 1$ bids p_K .
- Since the prices are market clearing, bidder k prefers buying position k than any other position m at these prices.

Vickrey Auction

- Bidders submit bids (\$ per-click)
- Seller finds assignment that maximizes total value assuming truthful bids
 - ▶ Puts highest bidder in top position, next in 2nd slot, etc.
- Charges each winner the total value their bid displaces.
 - ▶ For bidder n , each bidder below n is displaced by one position, so must add up the value of all these “lost” clicks.
- Note that Facebook uses a Vickrey auction.
- Dominant strategy to bid truthfully

Vickrey Auction Prices

- Bids: $b_1 > b_2 > \dots > b_N$.
- Bidder k displaces bidders $k + 1, \dots, K + 1$.
- Without k , each bidder $k + 1 < l < K + 1$ would get $x_{l-1}b_l$ instead of $x_l b_l$.
- Hence k should pay :

$$x_k p_k^V = \sum_{l=k+1}^{K+1} b_l (x_{l-1} - x_l) \quad \text{where } x_{K+1} = 0$$

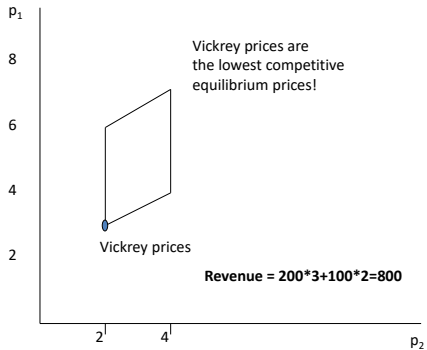
Theorem. Vickrey prices = lowest market clearing prices.

Vickrey Auction: Example

	Top (200 cpd)	2nd (100 cpd)
Bidder 1 (10 pcv)	2000	1000
Bidder 2 (4 pcv)	800	400
Bidder 3 (2 pcv)	400	200

- **Truthful bids:** 10,4,2
- Bidder 2 pays 200 for displacing bidder 3 (= \$2 per click).
- Bidder 1 pays 200 for displacing bidder 3 + 800-400 for displacing bidder 2 to slot 2 =600 (\$3 per click)

Vickrey Auction: Example



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Summary of Results

GSP Auction:

- Truthful bidding not a dominant strategy or an equilibrium in general.
- All market clearing prices are equilibrium prices (equivalent if we rule out strange equilibria).
- Edelman, Ostrovsky and Schwartz (2007, AER) argue that dynamics converge to lowest market clearing prices

Vickrey Auction:

- Truthful bidding is a dominant strategy.
- Vickrey prices are the lowest market clearing prices.
- Price computation is harder to understand for bidders.

Keyword Auction Design

- Platforms do retain some control over prices
 - ▶ Restricting the number of slots can increase prices.
 - ▶ Setting a reserve price can increase prices
- Platforms can also “quality-adjust” bids
 - ▶ In practice, ads that are more “clickable” get promoted.
 - ▶ Bids can be ranked according to $\text{score} = \text{bid} \times \text{quality}$.
 - ▶ This gives an advantage to high-quality advertisements.

Quality Scoring

- Suppose that instead of any bidder getting x_k clicks in position k , bidder n can expect to get $a_n x_k$ clicks.
- If a bidder has a high a_n , its ad is “clickable”.
- In practice, Google and Bing run giant regressions to try to estimate the “clickability” of different ads.
- Then bids in the auction can be ranked by $a_n b_n$, which means that clickable ads get prioritized in the rankings.
- This can have advantages and disadvantages
 - ▶ Puts weight on what users want and rewards higher quality ads.
 - ▶ Sometimes can reduce revenue if one ad gets lots of clicks.
- Also need to run tests to learn about “clickability”.

Sponsored vs Organic Results

- Google and Bing show “organic” search results and “sponsored” in different places.
 - ▶ The assignment of positions on the page is different
 - ▶ Organic search results: use algorithm to assess “relevance”
 - ▶ Sponsored search results: use bids to assess “value”
- To some extent there is competition
 - ▶ If a site gets a good organic position, should it pay for another?
 - ▶ Search engines have to think about maximizing user experience but also about capturing revenue from advertisers.

Summary: Sponsored Search

- Search auctions create a real-time market in which advertising opportunities are allocated to bidders.
- Auction theory suggests why the “second-price” rules used in practice might be reasonably efficient.
 - ▶ GSP does not induce “truthful” bidding but it has efficient Nash equilibria with competitive prices.
 - ▶ Vickrey auction does induce truthful bidding, but prices depend on a more complicated formula

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